# PARALLEL SUBGRADIENT METHOD FOR NONSMOOTH CONVEX OPTIMIZATION WITH A SIMPLE CONSTRAINT 

KAZUHIRO HISHINUMA AND HIDEAKI IIDUKA


#### Abstract

In this paper, we consider the problem of minimizing the sum of nondifferentiable, convex functions over a closed convex set in a real Hilbert space, which is simple in the sense that the projection onto it can be easily calculated. We present a parallel subgradient method for solving it and the two convergence analyses of the method. One analysis shows that the parallel method with a small constant step size approximates a solution to the problem. The other analysis indicates that the parallel method with a diminishing step size converges to a solution to the problem in the sense of the weak topology of the Hilbert space. Finally, we numerically compare our method with the existing method and state future work on parallel subgradient methods.


## 1. Introduction

This paper considers the following standard nonsmooth convex minimization problem.

Problem 1.1. Let $f_{i}(i=1,2, \ldots, K)$ be convex, continuous functionals on a real Hilbert space $H$ and let $C$ be a nonempty, closed convex subset of $H$. Then,

$$
\operatorname{minimize} \sum_{i=1}^{K} f_{i}(x) \text { subject to } x \in C
$$

A useful algorithm for solving Problem 1.1 is the incremental subgradient method [8, 12], and it is defined as follows: for defining $P_{C}$ as the projection onto $C$ and $\partial f_{i}(x)$ as the subdifferential of $f_{i}$ at $x \in H(i=1,2, \ldots, K)$, an iteration $(n+1)$ of the algorithm is

$$
\left\{\begin{array}{l}
\psi_{0, n}:=x_{n}  \tag{1.1}\\
\psi_{i, n}:=P_{C}\left(\psi_{i-1, n}-\lambda_{n} g_{i, n}\right), g_{i, n} \in \partial f_{i}\left(\psi_{i-1, n}\right) \quad(i=1,2, \ldots, K) \\
x_{n+1}:=\psi_{K, n}
\end{array}\right.
$$

[^0]Algorithm (1.1) requires us to use $P_{C}$ each iteration. Hence, we assume that $C$ is simple in the sense that $P_{C}$ can be easily calculated within a finite number of arithmetic operations [1, p.406], [2, Subchapter 28.3]. Some incremental methods that can be applied when $C$ is not always simple were presented in $[4,5,6]$.

Meanwhile, parallel proximal algorithms [2, Proposition 27.8], [3, Algorithm 10.27], [10] are also useful for solving Problem 1.1. They use the proximity operator of a nondifferentiable, convex $f_{i}$ which maps every $x \in H$ to the unique minimizer of $f_{i}+(1 / 2)\|x-\cdot\|^{2}$, where $\|\cdot\|$ stands for the norm of $H$. The parallel gradient algorithms presented in $[5,6]$ work only when $f_{i}$ is differentiable and convex, and $C$ is not always simple.

This paper presents a parallel subgradient method for solving Problem 1.1. The proposed method does not use any proximity operators, in contrast to the algorithms in [2, Proposition 27.8], [3, Algorithm 10.27], [10]. Next, we present convergence analyses for the two step-size rules: a constant step-size rule and a diminishing step-size rule. We show that the proposed method with a small constant step size approximates a solution to Problem 1.1. We also show that the algorithm with a diminishing step size weakly converges to a solution to Problem 1.1.

This paper is organized as follows. Section 2 gives the mathematical preliminaries. Section 3 presents the parallel algorithm for minimizing the sum of convex functionals over a simple, convex closed constraint set and studies its convergence properties for a constant step size and a diminishing step size. Section 4 provides numerical examples of the algorithm. Section 5 concludes the paper and mentions future work on parallel subgradient methods.

## 2. Preliminaries

Let $H$ be a real Hilbert space with inner product $\langle\cdot, \cdot\rangle$, and its induced norm $\|\cdot\|$. Let $\mathbb{N}$ denote the set of all positive integers including zero.
2.1. Subdifferentiability and projection. The subdifferential [2, Definition 16.1], [11, Section 23], [13, p.132] of $f: H \rightarrow \mathbb{R}$ is the set-valued operator,

$$
\partial f: H \rightarrow 2^{H}: x \mapsto\{u \in H: f(y) \geq f(x)+\langle y-x, u\rangle \quad(y \in H)\}
$$

Suppose that $f: H \rightarrow \mathbb{R}$ is continuous and convex with $\operatorname{dom}(f):=\{x \in H: f(x)<$ $\infty\}=H$. Then, $\partial f(x) \neq \emptyset(x \in H)$ [2, Proposition 16.14(ii)].

Proposition 2.1. [2, Proposition $16.14(\mathrm{iii})]$ Let $f: H \rightarrow \mathbb{R}$ be continuous and convex with $\operatorname{dom}(f)=H$. Then, for all $x \in H$, there exists $\delta>0$ such that $\partial f(B(x ; \delta))$ is bounded, where $B(x ; \delta)$ stands for a closed ball with center $x$ and radius $\delta$.

The metric projection [2, Subchapter 4.2, Chapter 28] onto a nonempty, closed convex set $C(\subset H)$ is denoted by $P_{C}$. It is defined by $P_{C}(x) \in C$ and $\| x-$ $P_{C}(x)\left\|=\inf _{y \in C}\right\| x-y \|(x \in H) . P_{C}$ is (firmly) nonexpansive with $\operatorname{Fix}\left(P_{C}\right):=$ $\left\{x \in H: P_{C}(x)=x\right\}=C[2$, Proposition 4.8, (4.8)].
2.2. Main problem. This paper deals with a networked system with $K$ users. Throughout this paper, we assume the following.

## Assumption 2.2.

(A1) $C(\subset H)$ is a nonempty, closed convex set, and $P_{C}$ can be easily calculated;
(A2) $f_{i}: H \rightarrow \mathbb{R}(i=1,2, \ldots, K)$ is continuous and convex with $\operatorname{dom}\left(f_{i}\right)=$ $\operatorname{dom}\left(\partial f_{i}\right)=H$
(A3) User $i(i=1,2, \ldots, K)$ can use $P_{C}$ and $\partial f_{i}$;
(A4) User $i(i=1,2, \ldots, K)$ can communicate with all users.
The main objective of this paper is to solve the following problem.
Problem 2.3. Under Assumption 2.2, find a minimizer of $\sum_{i=1}^{K} f_{i}$ over $C$.
We will use the following propositions to prove one of our main theorems.
Proposition 2.4. [14, Lemma 1] Suppose that $\left\{a_{n}\right\}$ and $\left\{b_{n}\right\}$ are two sequences of nonnegative members such that $a_{n+1} \leq a_{n}+b_{n}$ for all $n \in \mathbb{N}$. If $\sum_{n=0}^{\infty} b_{n}<\infty$, then $\lim _{n \rightarrow \infty} a_{n}$ exists.

Proposition 2.5. [9, Lemma 1] Suppose that $\left\{x_{n}\right\} \subset H$ converges weakly to $x \in H$ and $y \neq x$. Then, $\liminf _{n \rightarrow \infty}\left\|x_{n}-x\right\|<\liminf _{n \rightarrow \infty}\left\|x_{n}-y\right\|$.

## 3. Parallel algorithm

We present a parallel algorithm for solving Problem 2.3.

## Algorithm 3.1.

Step 0 . All users set $x_{0} \in H$ arbitrarily and $\left\{\lambda_{n}\right\} \subset(0, \infty)$.
Step 1. User $i(i=1,2, \ldots, K)$ computes $y_{i, n} \in H$ as follows:

$$
\left\{\begin{array}{l}
g_{i, n} \in \partial f_{i}\left(x_{n}\right) \\
y_{i, n}:=P_{C}\left(x_{n}-\lambda_{n} g_{i, n}\right)
\end{array}\right.
$$

Step 2. User $i(i=1,2, \ldots, K)$ shares $y_{i, n}$ in Step 1 with all users and calculates $x_{n+1} \in H$ as follows:

$$
x_{n+1}:=\frac{1}{K} \sum_{i=1}^{K} y_{i, n}
$$

Step 3. Put $n:=n+1$, and go to Step 1 .
Assumption (A2) ensures that $\partial f_{i}\left(x_{n}\right) \neq \emptyset(i=1,2, \ldots, K, n \in \mathbb{N})[2$, Proposition 16.14(ii)]. Assumption (A3) implies that user $i(i=1,2, \ldots, K)$ can compute $y_{i, n}$. Moreover, (A4) guarantees that all users can calculate $x_{n}$ in Step 2.

The convergence analyses of Algorithm 3.1 depend on the following assumption.
Assumption 3.2. For $i=1,2, \ldots, K$, there exists $M_{i} \in \mathbb{R}$ such that

$$
\sup \left\{\|g\|: g \in \partial f_{i}\left(x_{n}\right), n \in \mathbb{N}\right\}<M_{i}
$$

Suppose that $C$ is bounded (e.g., $C$ is a closed ball). From $\left\{y_{i, n}\right\} \subset C(i=$ $1,2, \ldots, K),\left\{y_{i, n}\right\}(i=1,2, \ldots, K)$ is bounded. Accordingly, $\left\{x_{n}\right\}$ is bounded. Hence, (A2) and Proposition 2.1 ensure that Assumption 3.2 holds. Moreover, since (A1) and (A2) imply that $C \cap \operatorname{dom}(f)=C \neq \emptyset$ and $C$ is bounded, (A2) (the continuity and convexity of $f$ ) guarantees that Problem 2.3 has a solution [2, Proposition 11.14].

This paper uses the notation,

$$
\begin{aligned}
M & :=\max \left\{M_{i}: i=1,2, \ldots, K\right\} \\
f & :=\sum_{i=1}^{K} f_{i}, X:=\left\{x \in C: f(x)=\inf _{y \in C} f(y)\right\} .
\end{aligned}
$$

We give the following lemma to analyze the convergence of Algorithm 3.1.
Lemma 3.3. Suppose that Assumption 3.2 holds and $\left\{x_{n}\right\} \subset H$ is the sequence generated by Algorithm 3.1. Then, for any $y \in C$ and for any $n \in \mathbb{N}$, we have

$$
\left\|x_{n+1}-y\right\|^{2} \leq\left\|x_{n}-y\right\|^{2}-\frac{2 \lambda_{n}}{K}\left(f\left(x_{n}\right)-f(y)\right)+\lambda_{n}^{2} M^{2}
$$

Proof. Choose $n \in \mathbb{N}$ arbitrarily. The convexity of $\|\cdot\|^{2}$ and the nonexpansivity of $P_{C}$ with $\operatorname{Fix}\left(P_{C}\right)=C$ imply that, for all $y \in C$,

$$
\begin{aligned}
\left\|x_{n+1}-y\right\|^{2} & =\left\|\frac{1}{K} \sum_{i=1}^{K} P_{C}\left(x_{n}-\lambda_{n} g_{i, n}\right)-P_{C}(y)\right\|^{2} \\
& \leq \frac{1}{K} \sum_{i=1}^{K}\left\|P_{C}\left(x_{n}-\lambda_{n} g_{i, n}\right)-P_{C}(y)\right\|^{2} \\
& \leq \frac{1}{K} \sum_{i=1}^{K}\left\|\left(x_{n}-y\right)-\lambda_{n} g_{i, n}\right\|^{2}
\end{aligned}
$$

which, together with $\|x-y\|^{2}=\|x\|^{2}-2\langle x, y\rangle+\|y\|^{2}(x, y \in H)$, means that

$$
\begin{aligned}
\left\|x_{n+1}-y\right\|^{2} & \leq \frac{1}{K} \sum_{i=1}^{K}\left(\left\|x_{n}-y\right\|^{2}-2\left\langle x_{n}-y, \lambda_{n} g_{i, n}\right\rangle+\left\|\lambda_{n} g_{i, n}\right\|^{2}\right) \\
& =\left\|x_{n}-y\right\|^{2}-\frac{2 \lambda_{n}}{K} \sum_{i=1}^{K}\left\langle x_{n}-y, g_{i, n}\right\rangle+\frac{\lambda_{n}^{2}}{K} \sum_{i=1}^{K}\left\|g_{i, n}\right\|^{2}
\end{aligned}
$$

From the definition of $\partial f_{i}(x)(x \in H)$, Assumption 3.2, and $f:=\sum_{i=1}^{K} f_{i}$, we find that, for all $y \in C$,

$$
\begin{aligned}
\left\|x_{n+1}-y\right\|^{2} & \leq\left\|x_{n}-y\right\|^{2}-\frac{2 \lambda_{n}}{K} \sum_{i=1}^{K}\left(f_{i}\left(x_{n}\right)-f_{i}(y)\right)+\lambda_{n}^{2} M^{2} \\
& =\left\|x_{n}-y\right\|^{2}-\frac{2 \lambda_{n}}{K}\left(f\left(x_{n}\right)-f(y)\right)+\lambda_{n}^{2} M^{2}
\end{aligned}
$$

This completes the proof.
3.1. Constant step-size rule. In this subsection, we study the convergence of Algorithm 3.1 when the step size is some constant.

Theorem 3.4. Suppose that Assumption 3.2 holds. Let $\lambda$ be a positive real number and let $\left\{x_{n}\right\} \subset H$ be the sequence generated by Algorithm 3.1. When $\lambda_{n}:=\lambda$ for all $n \in \mathbb{N}$, the following holds.

$$
\liminf _{n \rightarrow \infty} f\left(x_{n}\right) \leq \inf _{x \in C} f(x)+\frac{1}{2} \lambda K M^{2}
$$

Proof. Assume that the assertion does not hold. There exists a positive real number $\epsilon_{1}$ which satisfies the following inequality:

$$
\inf _{x \in C} f(x)+\frac{1}{2} \lambda K M^{2}+\epsilon_{1} \leq \liminf _{n \rightarrow \infty} f\left(x_{n}\right)
$$

Choose a positive real number $\epsilon_{2}$ such that $\epsilon_{2}<\epsilon_{1}$. From the property of the lower bound of $f$ over $C$, there exists $y \in C$ such that

$$
f(y)<\inf _{x \in C} f(x)+\left(\epsilon_{1}-\epsilon_{2}\right)
$$

Hence, we have

$$
f(y)+\frac{1}{2} \lambda K M^{2}+\epsilon_{2}<\liminf _{n \rightarrow \infty} f\left(x_{n}\right)
$$

Let $\epsilon_{3}$ be a positive real number which satisfies $\epsilon_{3}<\epsilon_{2}$. The property of the limit inferior of $f$ guarantees that $k_{0} \in \mathbb{N}$ exists such that, for all $k \geq k_{0}$,

$$
\liminf _{n \rightarrow \infty} f\left(x_{n}\right)-\left(\epsilon_{2}-\epsilon_{3}\right) \leq f\left(x_{k}\right)
$$

Therefore, using the two preceding inequalities, we have that, for all $k \geq k_{0}$,

$$
\frac{1}{2} \lambda K M^{2}+\epsilon_{3}<f\left(x_{k}\right)-f(y)
$$

Therefore, Lemma 3.3 ensures that, for all $k \geq k_{0}$,

$$
\begin{aligned}
\left\|x_{k+1}-y\right\|^{2} & \leq\left\|x_{k}-y\right\|^{2}-\frac{2 \lambda}{K}\left(f\left(x_{k}\right)-f(y)\right)+\lambda^{2} M^{2} \\
& <\left\|x_{k}-y\right\|^{2}-\frac{2 \lambda}{K}\left(\frac{1}{2} \lambda K M^{2}+\epsilon_{3}\right)+\lambda^{2} M^{2} \\
& =\left\|x_{k}-y\right\|^{2}-\frac{2 \lambda \epsilon_{3}}{K}
\end{aligned}
$$

which implies that, for all $k>k_{0}$,

$$
0 \leq\left\|x_{k}-y\right\|^{2}<\left\|x_{k_{0}}-y\right\|^{2}-\frac{2 \lambda \epsilon_{3}}{K}\left(k-k_{0}\right)
$$

However, since there exists a natural number $k_{1}>k_{0}$ such that

$$
\left\|x_{k_{0}}-y\right\|^{2}<\frac{2 \lambda \epsilon_{3}}{K}\left(k_{1}-k_{0}\right)
$$

we arrive at a contradiction. Therefore, $\liminf _{n \rightarrow \infty} f\left(x_{n}\right) \leq \inf _{x \in C} f(x)+\left(\lambda K M^{2}\right) / 2$ holds. This completes the proof.
3.2. Diminishing step-size rule. The main objective of this subsection is to prove the sequence generated by Algorithm 3.1 converges weakly to some point of the solution set $X$ of Problem 2.3. We first show the following.
Lemma 3.5. Suppose that Assumption 3.2 holds and $\left\{x_{n}\right\} \subset H$ is the sequence generated by Algorithm 3.1, with $\left\{\lambda_{n}\right\}$ satisfying

$$
\lim _{n \rightarrow \infty} \lambda_{n}=0 \text { and } \sum_{n=0}^{\infty} \lambda_{n}=\infty
$$

If $X$ is nonempty, we have

$$
\liminf _{n \rightarrow \infty} f\left(x_{n}\right)=\min _{x \in C} f(x)
$$

Proof. Assume that the assertion does not hold; i.e., $\min _{x \in C} f(x)<\liminf _{n \rightarrow \infty} f\left(x_{n}\right)$. Then, there exists a positive real number $\epsilon_{1}$ such that

$$
\min _{x \in C} f(x)+\epsilon_{1} \leq \liminf _{n \rightarrow \infty} f\left(x_{n}\right)
$$

The nonempty condition of $X$ guarantees the existence of $\hat{y} \in X$ satisfying

$$
f(\hat{y})=\min _{x \in C} f(x) \leq \liminf _{n \rightarrow \infty} f\left(x_{n}\right)-\epsilon_{1}
$$

Take a positive real number $\epsilon_{2}$ with $\epsilon_{2}<\epsilon_{1}$. The property of the limit inferior guarantees that $k_{1} \in \mathbb{N}$ exists such that, for all $k \geq k_{1}$,

$$
\liminf _{n \rightarrow \infty} f\left(x_{n}\right)-\left(\epsilon_{1}-\epsilon_{2}\right) \leq f\left(x_{k}\right)
$$

Using the two preceding inequalities, we find that, for all $k \geq k_{1}, f(\hat{y}) \leq f\left(x_{k}\right)+$ $\left(\epsilon_{1}-\epsilon_{2}\right)-\epsilon_{1}$; i.e., for all $k \geq k_{1}$,

$$
\epsilon_{2} \leq f\left(x_{k}\right)-f(\hat{y})
$$

Lemma 3.3 ensures that, for all $k \geq k_{1}$,

$$
\begin{aligned}
\left\|x_{k+1}-\hat{y}\right\|^{2} & \leq\left\|x_{k}-\hat{y}\right\|^{2}-\frac{2 \lambda_{k}}{K}\left(f\left(x_{k}\right)-f(\hat{y})\right)+\lambda_{k}^{2} M^{2} \\
& \leq\left\|x_{k}-\hat{y}\right\|^{2}-\frac{2 \lambda_{k}}{K} \epsilon_{2}+\lambda_{k}^{2} M^{2} \\
& =\left\|x_{k}-\hat{y}\right\|^{2}-\lambda_{k}\left(\frac{2}{K} \epsilon_{2}-\lambda_{k} M^{2}\right)
\end{aligned}
$$

Choose a positive real number $\epsilon_{3}$ such that $\epsilon_{3}<(2 / K) \epsilon_{2}$. The convergence of $\left\{\lambda_{n}\right\}$ to 0 implies the existence of $k_{2} \in \mathbb{N}$ such that, for all $k \geq k_{2}$,

$$
\lambda_{k}<\frac{1}{M^{2}}\left(\frac{2}{K} \epsilon_{2}-\epsilon_{3}\right)
$$

Therefore, putting $k_{3}:=\max \left\{k_{1}, k_{2}\right\}$, we have that, for all $k \geq k_{3}$,

$$
\left\|x_{k+1}-\hat{y}\right\|^{2}<\left\|x_{k}-\hat{y}\right\|^{2}-\lambda_{k} \epsilon_{3}
$$

which implies that, for all $k>k_{3}$,

$$
\begin{equation*}
\left\|x_{k}-\hat{y}\right\|^{2}<\left\|x_{k_{3}}-\hat{y}\right\|^{2}-\epsilon_{3} \sum_{n=k_{3}}^{k-1} \lambda_{n} \tag{3.1}
\end{equation*}
$$

The condition $\sum_{n=0}^{\infty} \lambda_{n}=\infty$ and (3.1) lead us to a contradiction. Therefore, $\liminf _{n \rightarrow \infty} f\left(x_{n}\right)=\min _{x \in C} f(x)$ holds. This completes the proof.

Now, we are in the position to perform the convergence analysis on Algorithm 3.1.

Theorem 3.6. Suppose that the assumptions in Lemma 3.5 hold and $\left\{x_{n}\right\} \subset H$ is the sequence generated by Algorithm 3.1, with $\left\{\lambda_{n}\right\}$ satisfying

$$
\sum_{n=0}^{\infty} \lambda_{n}=\infty \text { and } \sum_{n=0}^{\infty} \lambda_{n}^{2}<\infty
$$

Then, $\left\{x_{n}\right\}$ converges weakly to some point in $X$.
Proof. Lemma 3.5 guarantees that

$$
\liminf _{n \rightarrow \infty} f\left(x_{n}\right)=\min _{x \in C} f(x)
$$

which implies that there exists a subsequence $\left\{x_{n_{i}}\right\}$ of $\left\{x_{n}\right\}$ such that

$$
\lim _{i \rightarrow \infty} f\left(x_{n_{i}}\right)=\min _{x \in C} f(x)
$$

From the convexity of $C$ and the fact that $y_{i, n} \in C$ for any $i \in\{1,2, \ldots, K\}$ and for any $n \in \mathbb{N}$, Step 2 of Algorithm 3.1 guarantees that $x_{k} \in C$ for all natural numbers $k \geq 1$. Hence, $f\left(x^{\star}\right) \leq f\left(x_{k}\right)$ for any $x^{\star} \in X$ and for all $k \geq 1$. Therefore, Lemma 3.3 ensures that, for any $x^{\star} \in X$ and for all $k \geq 1$,

$$
\begin{align*}
\left\|x_{k+1}-x^{\star}\right\|^{2} & \leq\left\|x_{k}-x^{\star}\right\|^{2}-\frac{2 \lambda_{n}}{K}\left(f\left(x_{k}\right)-f\left(x^{\star}\right)\right)+\lambda_{k}^{2} M^{2} \\
& \leq\left\|x_{k}-x^{\star}\right\|^{2}+\lambda_{k}^{2} M^{2}  \tag{3.2}\\
& \leq\left\|x_{1}-x^{\star}\right\|^{2}+M^{2} \sum_{n=1}^{k} \lambda_{n}^{2}
\end{align*}
$$

Inequality (3.2) and $\sum_{n=0}^{\infty} \lambda_{n}^{2}<\infty$ mean that $\left\{x_{n}\right\}$ is bounded. The boundedness of $\left\{x_{n}\right\}$ and the closedness of $C$ guarantee the existence of a subsequence $\left\{x_{n_{i_{j}}}\right\}$ of $\left\{x_{n_{i}}\right\}$ that converges weakly to some point $z \in C$. Since (A2) implies that $f$ is continuous and convex, $f$ is weakly lower semicontinuous [2, Theorem 9.1]; i.e., $f(z) \leq \liminf _{j \rightarrow \infty} f\left(x_{n_{i_{j}}}\right)$. Therefore,

$$
\min _{x \in C} f(x) \leq f(z) \leq \liminf _{j \rightarrow \infty} f\left(x_{n_{i_{j}}}\right)=\lim _{j \rightarrow \infty} f\left(x_{n_{i_{j}}}\right)=\min _{x \in C} f(x)
$$

which implies that $z \in X$. Moreover, (3.2) and Proposition 2.4 lead us to the existence of $\lim _{n \rightarrow \infty}\left\|x_{n}-x^{\star}\right\|$ for all $x^{\star} \in X$. Let us take another subsequence $\left\{x_{n_{i_{k}}}\right\}$ of $\left\{x_{n_{i}}\right\}$ such that $\left\{x_{n_{i_{k}}}\right\}$ weakly converges to $w \in H$. A similar discussion to the one for obtaining $z \in X$ guarantees that $w \in X$. Here, we shall prove that $z=w$. Let us assume that $z \neq w$. Then, the existence of $\lim _{n \rightarrow \infty}\left\|x_{n}-x^{\star}\right\|$ for all $x^{\star} \in X$ and Opial's condition (Proposition 2.5) imply that

$$
\begin{aligned}
\lim _{n \rightarrow \infty}\left\|x_{n}-z\right\| & =\lim _{j \rightarrow \infty}\left\|x_{n_{i_{j}}}-z\right\|<\lim _{j \rightarrow \infty}\left\|x_{n_{i_{j}}}-w\right\|=\lim _{n \rightarrow \infty}\left\|x_{n}-w\right\| \\
& =\lim _{k \rightarrow \infty}\left\|x_{n_{i_{k}}}-w\right\|<\lim _{k \rightarrow \infty}\left\|x_{n_{i_{k}}}-z\right\|=\lim _{n \rightarrow \infty}\left\|x_{n}-z\right\|
\end{aligned}
$$

which is a contradiction. Accordingly, any subsequence of $\left\{x_{n_{i}}\right\}$ weakly converges to $z \in X$; i.e., $\left\{x_{n_{i}}\right\}$ weakly converges to $z \in X$. This implies that $z$ is a weak cluster point of $\left\{x_{n}\right\}$ and belongs to $X$. Moreover, since the existence of $\lim _{n \rightarrow \infty}\left\|x_{n}-x^{\star}\right\|$ for all $x^{\star} \in X$ guarantees that there is only one weak cluster point of $\left\{x_{n}\right\}$, the whole sequence $\left\{x_{n}\right\}$ weakly converges to $z \in X$. This completes the proof.

## 4. Numerical examples

We applied the incremental subgradient method (1.1) and Algorithm 3.1 to the following $N$-dimensional constrained nonsmooth convex optimization problem (Problem 1.1 when $H=\mathbb{R}^{N}$ and $K=N$ ).
Problem 4.1. Let $f_{i}: \mathbb{R}^{N} \rightarrow \mathbb{R}(i=1,2, \ldots, N)$ be convex and let $C$ be a nonempty, closed convex subset of $\mathbb{R}^{N}$. Then,

$$
\operatorname{minimize} \sum_{i=1}^{N} f_{i}(x) \text { subject to } x \in C \text {. }
$$

In the experiment, we used the PC-Cluster composed of 48 Fujitsu PRIMERGY RX350 S7 computers at the Ikuta campus of Meiji University. One of those computers has two Xeon E5-2690 (2.9GHz, 8 cores) CPUs and 32 GB memory. We used 64 CPU cores of this cluster; i.e., there were 64 users in the experiment environment that satisfied (A3) and (A4) of the Assumption 2.2. In the implementation of Step 2 in Algorithm 3.1, we used the MPI_Allreduce function, which is categorized as an All-To-All collective operation in [7, Chapter 5], to compute and share the sum of $y_{i, n}$ with all users. This means that all users contributed to computing $x_{n+1}$ in Algorithm 3.1. This operation does not violate Assumption 2.2. The experimental programs were written in C and compiled by gcc version 4.4 .7 with $\operatorname{Intel}(R)$ MPI Library 4.1. We used GNU Scientific Library 1.16 to express and compute vectors.

We set $N:=64$ and $C:=\left\{x \in \mathbb{R}^{N}:\|x\| \leq 1\right\}$ in Problem 4.1. For all $i=$ $1,2, \ldots, N$, we prepared random numbers $a_{i} \in(0,1)$ and $b_{i} \in(-1,1)$ and gave $a_{i}$ and $b_{i}$ to user $i$ in advance. The objective function of user $i$ was defined for all $x \in \mathbb{R}^{N}$ by $f_{i}(x):=\left|a_{i}\left\langle x, e_{i}\right\rangle+b_{i}\right|$, where $e_{i}(i=1,2, \ldots, N)$ stands for the natural base of $\mathbb{R}^{N}$.

In the experiment, we set $\lambda_{n}:=1$ for the constant step-size rule and $\lambda_{n}:=$ $1 /(n+1)$ for the diminishing step-size rule. We performed 100 samplings, each starting from the different random initial points in $[0,1)^{N}$.

Figure 1 shows the behaviors of $f(x):=\sum_{i=1}^{N} f_{i}(x)$ for the incremental subgradient method (1.1) and Algorithm 3.1 with a constant step size. The y-axes in Figures 1(a) and 1(b) represent the value of $f(x)$. The x -axis in Figure 1(a) represents the number of iterations and the x -axis in Figure 1(b) represents the elapsed time. The results show that Algorithm 3.1 minimizes the value of $f(x)$ more than the incremental subgradient method does (1.1).

Figure 2 shows the behaviors of $f(x)$ for the incremental subgradient method (1.1) and Algorithm 3.1 with the diminishing step size. The y-axes in Figures 2(a) and 2(b) represent the value of $f(x)$. The x-axis in Figure 2(a) represents the number of iterations, and the x-axis in Figure 2(b) represents the elapsed time. The results show that Algorithm 3.1 converges slower than the incremental subgradient


Figure 1. Behavior of $f(x)$ for the incremental subgradient method and Algorithm 3.1 with constant step size


Figure 2. Behavior of $f(x)$ with diminishing step size
method. However, it shows that Algorithm 3.1 with a constant step size behaves roughly to the same as the incremental subgradient method with the diminishing step size. This implies that, if it is difficult to share the diminishing step size with all users, Algorithm 3.1 can be used as an effective approximation algorithm of the incremental subgradient method.

## 5. Conclusion and future work

This paper discussed the problem of minimizing the sum of nondifferentiable, convex functions over a simple convex closed constraint set of a real Hilbert space. It presented a parallel algorithm for solving the problem. We studied its convergence properties for a constant step size and a diminishing step size. We showed that the algorithm with a constant step size approximates a solution to the problem, while the algorithm with a diminishing step size weakly converges to a solution to the problem. Finally, we numerically compared the algorithm with the existing algorithm and showed that, when the step size is constant, the algorithm performs better than the existing algorithm.

The numerical comparisons also indicated that, when the step size is diminishing, the existing algorithm converges to a solution faster than our algorithm. Therefore, in the future, we should consider developing parallel optimization algorithms which perform better than the existing algorithm even when the step sizes are diminishing.

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K. Hishinuma

Department of Computer Science, Meiji University, 1-1-1 Higashimita, Tama-ku, Kawasaki-shi, Kanagawa 214-8571, Japan

E-mail address: kaz@cs.meiji.ac.jp
H. Iiduka

Department of Computer Science, Meiji University, 1-1-1 Higashimita, Tama-ku, Kawasaki-shi, Kanagawa 214-8571, Japan

E-mail address: iiduka@cs.meiji.ac.jp


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